

What is claimed is:

1. A method for use in a public-key encryption system, the encryption system having an encryption block encrypting a plaintext  $m$  of a length of  $k_0$  to output a ciphertext  $(\alpha, \beta)$  and a decryption block for decrypting the ciphertext  $(\alpha, \beta)$  to provide the plaintext  $m$ , comprising the steps of:

(a) choosing variables  $p$ ,  $q$  and  $g$  as public-key parameters, wherein  $p$  is a large prime number of length  $k$ ,  $q$  is a large prime number dividing  $p-1$  and  $g$  is a generator for a multiplicative group  $Z_p^*$ , wherein  $Z_p^* = \{g^0, g^1, g^2, \dots, g^{q-1}\}$ ;

(b) choosing and publishing a first hash function  $H$ ,  $H: \{0, 1\}^k \rightarrow Z_q$ , providing security against an adaptive-chosen-ciphertext-attack and a second hash function  $G$ ,  $G: Z_p^* \rightarrow \{0, 1\}^k$ , providing security under a computational Diffie-Hellman assumption;

(c) choosing and storing a secret key  $x$  satisfying  $x \in Z_q$  based on the chosen public-key parameters  $p$ ,  $q$  and  $g$  and generating a public key  $X$  ( $X = g^x$ ), thereby publishing the public-key parameters  $p$ ,  $q$  and  $g$  and the public key  $X$ ;

(d) encrypting the plaintext  $m$  by using the public key  $X$ , thereby generating the ciphertext  $(\alpha, \beta)$ ;

(e) verifying whether the ciphertext  $(\alpha, \beta)$  is valid or not; and

(f) if the ciphertext  $(\alpha, \beta)$  is verified to be valid, decrypting the ciphertext  $(\alpha, \beta)$  by using the secret key  $x$  to

recover the plaintext  $m$ .

2. The method of claim 1, wherein the ciphertext  $(\alpha, \beta)$  is defined as:

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$$(\alpha, \beta) = (g^{H(m\|r)}, G(X^{H(m\|r)} \bmod p) \oplus (m\|r))$$

where  $r$  is a random string of a length  $k_1$  with  $k_0 + k_1 = k$ .

10 3. The method of claim 2, wherein the verifying step (e) includes the step of (e1) computing  $t = G(\alpha^x) \oplus \beta$  and determining whether  $\alpha$  of the ciphertext  $(\alpha, \beta)$  is identical to  $g^{H(i)}$  or not.

15 4. The method of claim 3, wherein the decrypting step (f) includes the step of removing the random number  $r$  from  $t$  to thereby recover the plaintext  $m$ .

5. The method of claim 2, wherein the exponentiation  
20 operation is replaced by addition operation over elliptic curve group.